

Last class.

Theorem: Assume $\gcd(s, t) = 1$

$$\Rightarrow U(st) \cong U(s) \oplus U(t)$$

Proof. already showed last time:

$$\text{if } \underline{\Phi}: x \in U(st) \mapsto (x \bmod s, x \bmod t)$$

$$\text{then } \underline{\Phi}(xy) = \underline{\Phi}(x)\underline{\Phi}(y)$$

still to show:

1-1 and onto:

$$\text{assume } \underline{\Phi}(x) = \underline{\Phi}(y)$$

$$\Rightarrow x \bmod s = y \bmod s$$

$$x \bmod t = y \bmod t$$

$$\left. \begin{array}{l} \Rightarrow s \mid (x-y) \\ t \mid (x-y) \end{array} \right\}$$

$$\text{lcm}(s, t) \mid (x-y) \leftarrow$$

$$\text{But } \text{lcm}(s, t) = \frac{st}{\text{gcd}(s, t)} = st$$

$$\Rightarrow st \mid x - y \iff x \bmod st = y \bmod st$$

i.e. x and y same in $U(st)$

onto follows from Φ being $\mathbb{1}$ and from

$$|\Phi(U(st))| = |U(st)| = \phi(st) = \phi(s)\phi(t) =$$

\uparrow Euler ϕ function \nwarrow fact from number theory
 $= \#\{0 < j < st, \text{gcd}(j, st) = 1\}$

Φ is \mathbb{H}

$$= |U(s)| |U(t)| = |U(s) \oplus U(t)|$$

As $\Phi(U(st)) \subset U(s) \oplus U(t) \Rightarrow$ equality.

Example: $U(15) \cong U(3) \oplus U(5)$
 $\cong \mathbb{Z}_2 \oplus \mathbb{Z}_4$

Remark: ① One can show

$$U(p^n) \cong \mathbb{Z}_{(p-1)p^{n-1}} \quad \text{cyclic for } p > 2$$

$$U(2^n) \cong \mathbb{Z}_2 \oplus \mathbb{Z}_{2^{n-2}}$$

Can use this to determine structure of $U(n)$ for any n .

e.g. $U(120) = U(2^3 \cdot 3 \cdot 5)$
 $\cong U(2^3) \oplus U(3) \oplus U(5)$
 $\cong \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_4$

② If $n = pq$ p, q primes

$$\Rightarrow U(n) \cong U(p) \oplus U(q) \\ \cong \mathbb{Z}_{p-1} \oplus \mathbb{Z}_{q-1}$$

in particular $|U(n)| = (p-1)(q-1)$

If we know prime factorization $n = pq$

\Rightarrow easy to calculate $|U(n)| = (p-1)(q-1)$

\Rightarrow " " " $j^m \pmod n$ for large m

(because $j^{(p-1)(q-1)} = 1 \pmod n$)

If we do not know prime factorization $\forall j \in U(n)$

$j^m \pmod n = ?$ hard

key ingredient for RSA encryption (see book for details
Chapter 8, p. 165)

Chapter 10 Homomorphisms

Def. G, H groups

A map $\underline{\Phi}: G \rightarrow H$ is called a **homomorphism**

if $\underline{\Phi}(ab) = \underline{\Phi}(a) \underline{\Phi}(b)$ for all $a, b \in G$

Remark: unlike an isom., a homomorphism need not be 1-1 or onto.

Examples: ① Every isom. is a homom.

② $G = \text{Gel}(2, \mathbb{R})$, $H = \mathbb{R}^* = (\mathbb{R} \setminus \{0\}, \cdot)$
Then $\underline{\Phi}(A) = \det(A)$ for $A \in \text{Gel}(2, \mathbb{R})$ is a homom.

proof: $\underline{\Phi}(AB) = \det(AB) \stackrel{\text{lin. algebra}}{=} \det(A) \det(B) = \underline{\Phi}(A) \underline{\Phi}(B)$ ✓

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$$G = S_m \quad H = \mathbb{Z}_2$$

recall: A perm. π is called odd/even

if $\pi = s_1 s_2 \dots s_k$ where each s_i is a 2-cycle

such that k is odd/even.

Define

$$\Phi: S_m \rightarrow \mathbb{Z}_2$$

$$\pi \rightarrow \begin{cases} 0 & \text{if } \pi \text{ even} \\ 1 & \text{if } \pi \text{ odd} \end{cases}$$

check hom. property case by case

$$\text{let } \pi = s_1 \dots s_k, \quad \sigma = t_1 \dots t_\ell \quad s_i, t_j \text{ 2-cycles}$$

check for all 4 cases k, ℓ odd/even that hom. prop. satisfied.

$$\text{eg. } k \text{ odd, } \ell \text{ even} \Rightarrow k + \ell \text{ odd}$$

$$\Phi(\pi) = 1, \quad \Phi(\sigma) = 0$$

$$\pi\sigma = \underbrace{s_1 s_2 \dots s_k t_1 t_2 \dots t_\ell}_{k+\ell \text{ cycles}}$$

add. notation for \mathbb{Z}_2
odd perm. $\Rightarrow \Phi(\pi\sigma) = 1 = \Phi(\pi) + \Phi(\sigma)$

Properties of Homomorphisms (see book Th. 10.1)

Let $\underline{\Phi}: G \rightarrow H$ be a homom.

e_G and e_H identities of G and H

$$\textcircled{1} \quad \underline{\Phi}(e_G) = e_H$$

$$\textcircled{2} \quad \underline{\Phi}(g^n) = \underline{\Phi}(g)^n$$

$$\textcircled{2}' \quad \underline{\Phi}(g^{-1}) = \underline{\Phi}(g)^{-1}$$

(proof: $\underline{\Phi}(g^{-1}) \underline{\Phi}(g) \underset{\substack{\uparrow \\ \text{hom. prop}}}{=} \underline{\Phi}(g^{-1}g) = \underline{\Phi}(e_G) = e_H$)

simil. $\underline{\Phi}(g) \underline{\Phi}(g^{-1}) \underset{\substack{\uparrow \\ \text{hom. prop}}}{=} \dots = e_H$

by uniqueness of inverse $\underline{\Phi}(g^{-1})$ must be the inverse of $\underline{\Phi}(g)$

$$\textcircled{3} \quad \text{ord } \underline{\Phi}(g) \mid \text{ord}(g)$$

" $\underline{\Phi}(g)^{-1}$ "

proof. let $n = \text{ord}(g)$: $\Phi(g)^n = \Phi(g^n) = \Phi(e_G) = e_H$
 $\Rightarrow \text{ord } \Phi(g) \mid n$.

Remark: For homom. Φ $\text{ord } \Phi(g)$ can be smaller than $\text{ord}(g)$

Example: $\text{ord}(123) = 3$

but $\Phi: S_3 \rightarrow \mathcal{R}_2$ as in Ex. 3

we have $\Phi(123) = 0$ in \mathcal{R}_2
 \uparrow
even permutation

$\text{ord}(0) = 1 \neq \text{ord}(123) = 3$

Def. $\Phi: G \rightarrow H$ hom.

We define the kernel of Φ by

$$\ker \Phi = \{g \in G, \Phi(g) = e_H\}$$

Ex. $\Phi: S_n \rightarrow \mathbb{Z}_2$

$$\begin{aligned} \ker \Phi &= \{\pi \in S_n, \Phi(\pi) = 0 \pmod{2}\} \\ &= \{\pi \in S_n, \pi \text{ even}\} \\ &= A_n. \end{aligned}$$

Lemma (a) $\ker \Phi$ is a normal subgroup of G

$$(b) \quad \Phi(a) = \Phi(b) \iff a \ker \Phi = b \ker \Phi$$

Proof.

(a)

Subgroup test

$$g, k \in K \Rightarrow \Phi(g) = e_H = \Phi(k)$$

$$\Rightarrow \Phi(gk) = \Phi(g)\Phi(k) = e_H e_H = e_H$$

$$\Rightarrow gk \in \ker \Phi$$

$$g \in \ker \Phi: \Phi(g^{-1}) = \Phi(g)^{-1} = e_H^{-1} = e_H$$

$$\Rightarrow g^{-1} \in \ker \Phi$$